

Analysis And Implementation of Noekeon Algorithms For Encryption and Description of Text Data

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Encryption-decryption
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Abstract. As the times evolve, human needs increase. Including information needs. Therefore, sending and storing files through electronic media requires a process that can ensure the security and integrity of the file. To ensure the safety and integrity of a file, an encoding process is required. Encryption is performed when the file is sent. This process will convert the original file into a confidential unreadable file. Meanwhile, the decryption process is done by the recipient of the sent file. The personal file received will be converted back to the original file. By encoding, the original file will not be read by unauthorized parties but only by recipients who have a decryption key.

1. INTRODUCTION

Currently, one of the ways used for file security is to use a cryptographic system that is to encode the contents of the information (plaintext) into content that is not understood through the encryption process (chipper) and to regain the original data, the decryption process (decipher)—accompanied by using the correct key[1]. But with the development of the science of encoding or cryptography, efforts to obtain the key can be made by anyone, including those who are not authorized to have such information. Therefore, researchers on cryptography will constantly evolve to get increasingly powerful cryptographic algorithms, making efforts to break down cryptographic codes unlawfully become more complex[2].

Information communication today has expanded widely. Information communication in exchanging messages and transactions until the exchange of confidential information has used computer technology and a vast and complex global network. The breadth of information communication utilizing computer technology and global networks today led to the emergence of parties who want to steal such information[3]. The vast and complex global network makes it almost impossible to protect messages right to their destination. The threat of information theft has led many researchers to develop techniques or methods to provide security to the content of information in notes, one of which is data encoding techniques or cryptography[4].

Cryptography is the science and art of maintaining the confidentiality of messages by encoding them into a form that is no longer understood. The purpose of cryptography is to provide security in the form of / privacy, data integrity, authentication, non-repudiation[5]. In the field of cryptography, there are two main processes, encryption, and decryption[6]. Encryption is the process of encoding plaintext or messages into ciphertext. Decryption is the process of returning the ciphertext to its original plaintext[7].

2. METHOD

The early stage is the determination of the use of materials related to the research conducted. At this stage will be done identifying the problem and selecting algorithms and methods to solve it. Issues are found by following current technology issues and developments and studying research that has been done and published through scientific journals. The problem to be researched is how to secure data. To avoid theft or leakage of data on cloud storage, the encryption system is implemented with noekeonalgorithm. Stages with the following flow:

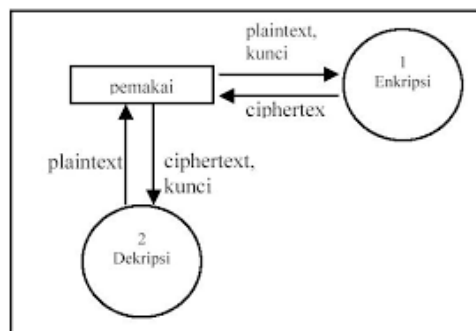


Figure 1. Noekeon Algorithm Work Process

2.1. Noekeon

Noekeon is a repeating code block with a block length and a key length of 128 bits each, consisting of a simple repeating spin transformation application, followed by an output transformation. Noekeon has 16 rounds (N_r) of iterations[8]. Each game is carried out three modifications, namely theta, offset shift consisting of two changes P_{i1} and P_{i2} , and gamma[9].

Critical scheduling is done by converting a 128-bit primary key (code-key) into a working key. In Noekeon, there is a mode in which key scheduling is not performed, called direct-key mode, which serves the key scheduling process to eliminate related-key attack patterns[10].

3. RESULT AND DISCUSSION

The encryption and decryption process of Noekeon's algorithm uses the key generation or scheduling process used in each round of encryption and decryption. The key stages of the neocon algorithm are as follows:

1. Input the user key.
2. They are padding the user key if the key length is less than 16 characters.
3. Form A0-A3 blocks of 32 bits or four characters each.
4. Performing a key scheduling round of 16 rounds (i), the operations on each game are as follows:
 1. Perform xor A0 operations with RC(i).
 2. Perform inverse theta operations against A0-A3.
 3. Perform rotate_left operations $A1 \ll 1$
 4. Perform operations rotate_left $A2 \ll 5$
 5. Perform operations rotate_left $A3 \ll 2$
 6. Perform gamma operations against A0-A3.
 7. Perform rotate_right operation $A1 \gg 1$
 8. Perform rotate_right operation $A2 \gg 5$
 9. Perform operations rotate_right $A3 \gg 2$
10. After the next round of operations, perform xor operations A0 and RC(16). 11. A0 – A3 will be the DK(0)-DK(3) decryption key.
12. Perform inverse theta operations against A0-A3.
13. A0-A3 will be the key to EK(0)-EK(3).

Table 1. RC Constants

i	Formula	RC(i)
0	-	0x80
1	$RC[0] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[1] = Rc[0] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[1]= Rc[0] \ll 1$	0x1a
2	$RC[1] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[2] = Rc[1] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[2]= Rc[1] \ll 1$	0x34
3	$RC[2] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[3] = Rc[2] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[3]= Rc[2] \ll 1$	0x68

4	$RC[3] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[4] = Rc[3] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[4]= Rc[3] \ll 1$	0xd0
5	$RC[4] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[5] = Rc[4] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[5]= Rc[4] \ll 1$	0xba
6	$RC[5] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[6] = Rc[5] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[6]= Rc[5] \ll 1$	0x6e
7	$RC[6] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[7] = Rc[6] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[7]= Rc[6] \ll 1$	0xdc
8	$RC[7] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[8] = Rc[7] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[8]= Rc[7] \ll 1$	0xa2
9	$RC[8] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[9] = Rc[8] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[9]= Rc[8] \ll 1$	0x5e
10	$RC[9] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[10] = Rc[9] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[10]= Rc[9] \ll 1$	0xbc
11	$RC[10] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[11] = Rc[10] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[11]= Rc[10] \ll 1$	0x62
12	$RC[11] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[12] = Rc[11] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[12]= Rc[11] \ll 1$	0xc4
13	$RC[12] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[13] = Rc[12] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[13]= Rc[12] \ll 1$	0x92
14	$RC[13] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[14] = Rc[13] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[14]= Rc[13] \ll 1$	0x3e
15	$RC[14] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[15] = Rc[14] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[15]= Rc[14] \ll 1$	0x7c
16	$RC[15] \text{ And } 0x80 \neq 0 ? \text{ If True } Rc[16] = Rc[15] \ll 1 \text{ Xor } 0x1B \text{ If False } RC[16]= Rc[15] \ll 1$	0xf8

Analysis of the key scheduling process in Noekeon's algorithm can be described as follows:

Key Scheduling :

Input Key : 123450000000000000

Key(Bit) :

00110001;00110010;00110011;00110100;

00110101;00110000;00110000;00110000;

00110000;00110000;00110000;00110000;

00110000;00110000;00110000;00110000;

A0 : 00110001;00110010;00110011;00110100

= 825373492

A1 : 00110101;00110000;00110000;00110000

= 892350512

A2 : 00110000;00110000;00110000;00110000

= 808464432

A3 : 00110000;00110000;00110000;00110000

= 808464432

Round (0)Scheduling Key :

A0 = A0 Xor RC(0) = 825373620

Theta Inv Process : A0 : 825373620

A1 : 892350512

A2 : 808464432

A3 : 808464432 T = A0 Xor A2

= 16909188

A0 = A0 Xor RC(16) = 2504158901 DK(0) =

A0 = 2504158901

DK(1) = A1 = 334622109 DK(2) = A2 =

3718075253 DK(3) = A3 = 205715165

Theta Inv Process :

A0 : 2504158901

A1 : 334622109

A2 : 3718075253

A3 : 205715165 T = A0 Xor A2

= 1222589888

Temp = T<<8 Xor T>>8 = 527769461 A1 = T

Xor Temp Xor A1 = 1146867496 A3 = T Xor

Temp Xor A3 = 1541985384 T = A1 Xor A3 =

531831616

Temp = T<<8 Xor T>>8 = 4077187844 A0 = T

Xor Temp Xor A0 = 2046134001 A2 = T Xor

Temp Xor A2 = 824884017

Temp = T<<8 Xor T>>8 = 2248312322 A1 = T

Xor Temp Xor A1 = 2989536694 A3 = T Xor

Temp Xor A3 = 3073422774 T = A1 Xor A3 =

83886080

Temp = T<<8 Xor T>>8 = 327685

A0 = T Xor Temp Xor A0 = 876032945 A2 = T

Xor Temp Xor A2 = 892678197 A1 = A1 << 1 =

1684106093

A2 = A2 << 5 = 2795898534 A3 = A3 << 2 =

3703756506

Gamma Process :

A0 : 876032945

A1 : 1684106093

A2 : 2795898534

A3 : 3703756506

A1 = A1 Xor ((Not A3) And (Not A2)) =

1702380140

A0 = A0 Xor (A2 And A1) = 269955477 A0 = A3

= 3703756506

A3 = A0 = 269955477

A2 = A2 Xor A3 Xor A1 Xor A0 = 252420997

A1 = A1 Xor (Not A3 And Not A2) =

2241334790

A0 = A0 Xor (A2 And A1) = 3653948638 A1 =

A1 >> 1 = 1120667395

A2 = A2 >> 5 = 678976796 A3 = A3 >> 2 =

1141230693

I(0) = A0 = 2046134001 I(1) = A1 = 1146867496

I(2) = A2 = 824884017 I(3) = A3 = 1541985384

Based on the analysis of key scheduling that has

been done, the key is obtained for the encryption

and decryption process as follows:

Input key : 12345 Encryption key :

I(0) = A0 = 2046134001 I(1) = A1 = 1146867496

I(2) = A2 = 824884017 I(3) = A3 = 1541985384

Decryption key :

DK(0) = A0 = 2504158901 DK(1) = A1 =

334622109 DK(2) = A2 = 3718075253 DK(3) =

A3 = 205715165

3.1 Encryption Process Analysis

Encryption analysis in Noekeon's algorithm uses input blocks of 32 bits in length. Here is the description of the operation of the encryption process in Noekeon's algorithm:

1. Input Plaintext from the user.
2. They are padding the plaintext if the mod length of 16 is not equal to 0.
3. Form A0-A3 input blocks that are 32 bits or four characters long, respectively.

4. Perform an encryption loop of 16 rounds(i) which can be described as follows.

1. Perform xor A0 operations with RC(i).
2. Perform theta operations against A0-A3 using K(0)-K(3).
3. Perform rotate_left operations $A1 \ll 1$
4. Perform operations rotate_left $A2 \ll 5$
5. Perform operations rotate_left $A3 \ll 2$
6. Perform gamma operations against A0-A3.
7. Perform rotate_right operation $A1 \gg 1$
8. Perform rotate_right operation $A2 \gg 5$
9. Perform operations rotate_right $A3 \gg 2$
10. After the next round of operations, perform xor operations A0 and RC(16).
11. Perform theta operations against A0-A3 using K(0)-K(3).
12. Break down blocks A0-A3 into bytes of ciphertext characters.

Analysis of the encryption process in Noekeon's algorithm can be described as follows.

Plaintext : "zizi encryption 12" Input Block	A0 : 1701735282
Formation :	A1 : 1768977257
Input Block (0) :	A2 : 544893306
01100101;01101110;01101011;01110010	A3 : 1763717426
Input Block (1) :	Round (0) Encryption :
01101001;01110000;01110011;01101001	A0 = A0 Xor RC(0) = 1701735410
Input Block (2) :	Theta Process : A0 : 1701735410
00100000;01111010;01101001;01111010	A1 : 1768977257
Input Block (3) :	A2 : 544893306
01101001;00100000;00110001;00110010	A3 : 1763717426
End (4) BLock encryption:	

$K(0) : 2046134001$
 $K(1) : 1146867496$
 $K(2) : 824884017$
 $K(3) : 1541985384$
 $T = A0 \text{ Xor } A2$
 $= 1158939272$
 $\text{Temp} = T \ll 8 \text{ Xor } T \gg 8 = 2621938759$
 $A1 = T \text{ Xor } \text{Temp} \text{ Xor } A1 = 2955144614$
 $A3 = T \text{ Xor } \text{Temp} \text{ Xor } A3 = 2960371709$
 $A0 = A0 \text{ Key} \text{ Xor}(0) = 479980803$
 $A1 = A1 \text{ Key} \text{ Xor}(1) = 4101513870$
 $A2 = A2 \text{ Key} \text{ Xor}(2) = 290509387$
 $A3 : 2926440023$
 $A1 = A1 \text{ Xor } ((\text{Not } A3) \text{ And } (\text{Not } A2)) =$
 3923792957
 $A0 = A0 \text{ Xor } (A2 \text{ And } A1) = 865995842$
 $A0 = A3 = 2926440023$
 $A3 = A0 = 865995842$
 $A2 = A2 \text{ Xor } A3 \text{ Xor } A1 \text{ Xor } A0 = 2864567542$
 Round operations are performed in 16 rounds. After the round is complete then the operation is continued as follows.
 $A0 = A0 \text{ Xor } \text{RC}(16) = 2594189566$
 Theta Process :
 $A0 : 2594189566$
 $A1 : 347116736$
 $A2 : 563119268$
 $A3 : 4101676072$
 $K(0) : 2046134001$
 $K(1) : 1146867496$
 $A2 \text{ Xor } \text{Key}(2) = 280641429$
 $A3 = A3 \text{ Xor } \text{Key}(3) = 2125305365$ $T = A1 \text{ Xor}$
 $A3 = 4286124968$
 $\text{Temp} = T \ll 8 \text{ Xor } T \gg 8 = 3521958380$ $A0 = T$
 $\text{Xor } \text{Temp} \text{ Xor } A0 = 3451941963$ $A2 = T \text{ Xor}$
 $\text{Temp} \text{ Xor } A2 = 1043332561$
 After the process is complete, the ciphertext can be obtained by breaking block A0-A3 into bytes of characters that can be seen as follows:
 $\text{Output}(0) : 205$
 $\text{Output}(1) : 192$
 $\text{Output}(2) : 116$
 $\text{Output}(3) : 75$
 $\text{Output}(4) : 129$
 $\text{Output}(5) : 212$

3.2 Decryption Process Analysis

The decryption process in Noekeon's algorithm uses input blocks of 32 bits in length. Here is the description of the operation of the

$A3 = A3 \text{ Xor } \text{Key}(3) = 3952835477$
 $T = A1 \text{ Xor } A3 = 534993179$
 $\text{Temp} = T \ll 8 \text{ Xor } T \gg 8 = 4165400646$
 $A0 = T \text{ Xor } \text{Temp} \text{ Xor } A0 = 4215163998$
 $A2 = T \text{ Xor } \text{Temp} \text{ Xor } A2 = 4143280918$
 $A1 = A1 \ll 1 = 3908060445$
 $A2 = A2 \ll 5 = 3735970526$
 $A3 = A3 \ll 2 = 2926440023$
 Gamma Process :
 $A0 : 4215163998$
 $A1 : 3908060445$
 $A2 : 3735970526$
 $A1 = A1 \text{ Xor } (\text{Not } A3 \text{ And } \text{Not } A2) =$
 2912974644
 $A0 = A0 \text{ Xor } (A2 \text{ And } A1) = 114145891$
 $A1 = A1 \gg 1 = 1456487322$
 $A2 = A2 \gg 5 = 3042307751$
 $A3 = A3 \gg 2 = 2363982608$

encryption process in Noekeon's algorithm:

- Input ciphertext from the user.
 $K(2) : 824884017$
 $K(3) : 1541985384$ $T = A0 \text{ Xor } A2$
 $= 3140531290$
 $\text{Temp} = T \ll 8 \text{ Xor } T \gg 8 = 1779395087$ $A1 =$
 $T \text{ Xor } \text{Temp} \text{ Xor } A1 = 3314502293$ $A3 = T$
 $\text{Xor } \text{Temp} \text{ Xor } A3 = 625295997$ $A0 = A0 \text{ Key}$
 $\text{Xor}(0) = 3814045199$
 $A1 = A1 \text{ Xor } \text{Key}(1) = 2178188733$
 $A2$ =
 $\text{Output}(6) : 133$
 $\text{Output}(7) : 189$
 $\text{Output}(8) : 62$
 $\text{Output}(9) : 47$
 $\text{Output}(10) : 253$
 $\text{Output}(11) : 209$
 $\text{Output}(12) : 126$
 $\text{Output}(13) : 173$
 $\text{Output}(14) : 150$
 $\text{Output}(15) : 21$

Then obtained the results of encryption as follows:

Plaintext : "Zizi encryption 12" Chiperteks :
 $^2\text{OsY}\text{¶}[\text{a}_\text{á}\text{ý}_\text{ê}\text{Û}\text{□}1\text{Š}_\text{(_}[\text{X}_\text{ä}^{\text{†}}\text{sr}\&\text{v}^{\text{''}}\text{½}\{\$

- Form A0-A3 input blocks that are 32 bits or 4 characters long, respectively.
- Decryption rounds of 16 rounds(i) are performed from rounds 16 – 1 (inverted) which can be described below.
- Perform theta operations against A0-A3 using K(0)-K(3).

5. Perform xor A0 operations with RC(i).
6. Perform rotate_left operations $A1 \ll 1$
7. Perform operations rotate_left $A2 \ll 5$
8. Perform operations rotate_left $A3 \ll 2$
9. Perform gamma operations against A0- A3.
10. Perform rotate_right operation $A1 \gg$
11. Perform rotate_right operation $A2 \gg$
12. Perform operations rotate_right $A3 \gg$
13. Perform theta operations against A0-A3 using K(0)-K(3).
14. Breaks block A0-A3 into bytes of ciphertext characters.

Analysis of the decryption process in Noekeon's algorithm can be described as follows: ciphertext:

2OsY¶a_áy_êÛŠ_(X_ä'sr&v'1/2{ Input
 Block Formation :

Input	Block	(0)	:
11001101;	11000000;	01110100;	01001011
Input	Block	(1)	:
10000001;	11010100;	10000101;	10111101
Input	Block	(2)	:
00111110;	00101111;	11111101;	11010001
Input	Block	(3)	:
01111110;	10101101;	10010110;	00010101

End (4) BLOK

A0 : 3451941963

A1 : 2178188733

A2 : 1043332561

A3 : 2125305365

Round (0) Decryption :

Theta Process :

A0 : 3451941963

A1 : 2178188733

A2 : 1043332561

A3 : 2125305365

K(0) : 2504158901

K(1) : 334622109

K(2) : 3718075253

K(3) : 205715165 T = A0 Xor A2

= 4092561818

Temp = $T \ll 8 \text{ Xor } T \gg 8 = 1970959738$ A1 = T

Xor Temp Xor A1 = 121731421 A3 = T Xor

Temp Xor A3 = 4164446965 A0 = A0 Key

Xor(0) = 1484921598

A1 = A1 Xor Key(1) = 347116736 A2 = A2 Key

Xor(2) = 3820138148 A3 = A3 Xor Key(3) =

4101676072 T = A1 Xor A3 = 3771336936

Temp = $T \ll 8 \text{ Xor } T \gg 8 = 585638632$

A0 = T Xor Temp Xor A0 = 2594189566 A2 =

T Xor Temp Xor A2 = 563119268 A0 = A0 Xor

RC(0) = 2594189318

A1 = $A1 \ll 1 = 694233472$ A2 = $A2 \ll 5 =$
 839947396 A3 = $A3 \ll 2 = 3521802403$

Gamma Process :

A0 : 2594189318

A1 : 694233472

A2 : 839947396

A3 : 3521802403

A1 = A1 Xor ((Not A3) And (Not A2)) =
 627319512

A0 = A0 Xor (A2 And A1) = 3131060358 A0 =

A3 = 3521802403

A3 = A0 = 3131060358

A2 = A2 Xor A3 Xor A1 Xor A0 = 2084501113

A1 = A1 Xor (Not A3 And Not A2) = 606415832

A0 = A0 Xor (A2 And A1) = 4123939579

A1 = $A1 \gg 1 = 303207916$ A2 = $A2 \gg 5 =$

3420583859 A3 = $A3 \gg 2 = 2930248737$

Round operations are performed in 16 rounds. After the round is complete then the operation is continued as follows.

Theta Process :

A0 : 4215163998

A1 : 4101513870

A2 : 4143280918

A3 : 3952835477

K(0) : 2504158901

K(1) : 334622109

K(2) : 3718075253

K(3) : 205715165 T = A0 Xor A2

= 231423816

Temp = $T \ll 8 \text{ Xor } T \gg 8 = 2201125682$ A1 = T

Xor Temp Xor A1 = 2055315188 A3 = T Xor

Temp Xor A3 = 1700972527 A0 = A0 Xor

Kunci(0) = 1853629163 A1 = A1 Xor Kunci(1) =

1768977257 A2 = A2 Xor Kunci(2) = 728244323

A3 = A3 Xor Kunci(3) = 1763717426 T = A1 Xor

A3 = 5259867

Temp = $T \ll 8 \text{ Xor } T \gg 8 = 188877634$ A0 = T

Xor Temp Xor A0 = 1701735410 A2 = T Xor

Temp Xor A2 = 544893306 A0 = A0 Xor RC(0)

= 1701735282

After the process is complete, the plaintext can be obtained by splitting block A0-A3 into byte characters that can be seen as follows:

Output(0) : 101
Output(1) : 110
Output(2) : 107
Output(3) : 114
Output(4) : 105
Output(5) : 112
Output(6) : 115
Output(7) : 105
Output(8) : 32

Output(9) : 122
Output(10) : 105
Output(11) : 122
Output(12) : 105
Output(13) : 32
Output(14) : 49
Output(15) : 50

Then obtained the results of decryption as follows:

Plaintext : "zizi encryption 12"

4. CONCLUSION

With the completion of this research which includes the process of encryption and decryption of messages using the Noekeon algorithm, it can be concluded that Noekeon's algorithm can be applied in encryption and decryption of text messages by dividing the characters contained in the text into blocks - Noekeon blocks so that it can be processed into noekeon rotation operations into ciphertext or plaintext. The process of designing and developing a text message encoding application using Noekeon's algorithm begins with analyzing Noekeon's algorithm both in the process and its provisions. The Next Stage is to create a neocon interface and process modules consisting of key initialization, key generation, encryption, and decryption.

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